



arm

# Writing Constant-Time Code

What, why and how?

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# Side-Channel Attacks

- Constant-time code is a countermeasure
- “Attack based on information gained (‘leaked’) from the practical implementation of a system”
  - Can completely break real world encryption
- Many types exist
  - Timing
  - Cache
  - Branch prediction
  - Many more (power, fault, etc...)

# A Simple Example

- Checking that a MAC (for example) matches an expected value:

```
int compare(uint8_t *a, uint8_t *b, size_t len)
{
    for (size_t i = 0; i < len; i++)
        if (a[i] != b[i])
            return -1;
    return 0;
}
```

- This is vulnerable to a number of side-channel attacks

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# What Can Be Done?

- Use constant-time code to eliminate this class of vulnerability
- Execution time cannot depend on secret values
- Follow the golden rules:
  - No branches depending on secret data
  - No memory access depending on secret data
  - No variable-time instruction executed on secret data
- Limits: physical side-channels, faults, readability

# Fixing Our Example

```
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# Fixing Our Example – Removing Branches

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    return 0;
}
```

# Fixing Our Example – Removing Branches

```
int compare(uint8_t *a, uint8_t *b, size_t len)
{
    char diff = 0;
    for (size_t i = 0; i < len; i++)
        diff |= a[i] ^ b[i];
    return diff;
}
```



# Example #2 - Table Lookup

- Seemingly a constant time instruction -  $O(1)$

```
u32 lookup(u32 *t, u32 l, u32 secret_i) {  
    return t[secret_i];  
}
```

- Vulnerable to a cache attack!
- Violates rule #2 (No memory access depending on secret data)

# Example #2 - Table Lookup

- Seemingly a constant time instruction -  $O(1)$

```
u32 lookup(u32 *t, u32 l, u32 secret_i) {
    u32 r = t[0];
    for (u32 j = 1; j < l; j++) {
        r = choose(r, t[j], eq(secret_i, j));
    }
    return r;
}
```

- Vulnerable to a cache attack!
- Violates rule #2 (No memory access depending on secret data)

# Real World Examples

- (1996) “Timing Attacks on Implementations of Diffie-Hellman, RSA, DSS, and Other Systems”
- (2003) “Remote Timing Attacks Are Practical”
- (2011) “Remote Timing Attacks Are Still Practical”
- (2013) Lucky 13
- (2018) Spectre & Meltdown

# Conclusion

- Constant-time code can be used to defend against a variety of side-channel attacks
- Constant-time code is not easy to write - prone to both human and compiler error
- Promising techniques to automate testing & writing
  - Still needs some improvements & widespread adoption

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Thank You

Danke

Merci

谢谢

ありがとう

Gracias

Kiitos

감사합니다

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